

REMARKS

Favorable reconsideration of this application is requested in view of the foregoing amendments and the following remarks. Claims 1-5 are pending in the application.

At page 2 of the Action, the Examiner objects to the drawings. Applicant submits herewith a proposed figure 2 that overcomes the objection to the drawings as set forth by the Examiner by showing the head and tail of the queues, the latest read pointer, committed read pointer, second latest read pointer, latest write pointer, and committed write pointer. Support for Figure 2 is found in table 2 at page 9, at page 14 last paragraph to page 15 third paragraph and in the claims as originally filed. The specification is also amended to recite a brief description of figure 2. Applicant requests that the Examiner approve figure 2 and withdraw the objection to the drawings. An annotated marked-up version of figure 2 is attached hereto.

Accordingly, withdrawal of this objection is respectfully requested.

At pages 2-3 of the Action, the Examiner objects to the claims. Claim 2 is amended as suggested by the Examiner. Claim 4 is amended as suggested by the Examiner.

Accordingly, withdrawal of this objection is respectfully requested.

Claim 5 stands rejected under 35 USC 112(2) as indefinite. The Examiner's careful consideration of, and various suggested changes to, the phraseology of the claims is appreciated. Claim 5 is amended to provide antecedent for the received data. Support for this change to claim 5 is found in the paragraph bridging pages 13-14 of this application as originally filed.

Accordingly, withdrawal of this rejection is respectfully requested.

Claims 1-5 stand rejected under 35 USC 102(b) as anticipated by Hamstra (i.e., U.S. Pat. No. 5,016,221). The disclosure of the Hamstra reference is not sufficient to support this rejection.

Hamstra fails to disclose a method of managing data stored in a queue wherein data is read from a queue and a latest read pointer is placed at the location corresponding to the end of the read data, the data is then transferred to a destination, and then, after receiving confirmation that the data transfer is successful, a committed read pointer is updated to a location in memory corresponding to the end of the data. With the invention of claim 1, therefore, the claimed "latest read pointer" always moves ahead of the "committed read pointer". As set out in the summary of the invention portion of the description of this application as originally filed, such an operation allows uncommitted data to be stored (for example for retransmission if required) without the provision of a separate area of memory.

Nowhere within Hamstra is such an arrangement comprising a committed read pointer following a latest read pointer disclosed. Instead, referring to Figures 3A to 3E of Hamstra, a read pointer 30, commit pointer 32, and write pointer 34 are provided for use with a FIFO RAM. As shown in Figure 3A of Hamstra, when the RAM is empty, the three pointers are all set at the same memory location (cf. Column 7 lines 13-15 of Hamstra). When information has begun to be received and stored within the RAM of Hamstra, the write pointer 34 is advanced, leaving a block 50 of stored, but as yet uncommitted, binary information within the RAM (cf. Column 7 lines 17 to 23, and Figure 3B of Hamstra). If it is decided that the stored block of data 50 is to be retained for use, the commit pointer 32 is advanced to coincide with the write pointer 34 (cf. Figure 3C, and column 7 lines 31 to 35). Committed data 50 between the read pointer 30 and the commit pointer 32 of Hamstra is now available to be read out from the RAM (cf. Column 7 lines 35-37). Such a read operation is illustrated in Figure 3E of Hamstra, wherein, as the committed binary information 50 is read out, the read pointer 30 advances. The read pointer 30 of Hamstra may advance up to, but no further than, the commit pointer 32 (cf. Column 8 lines 50 to 53).

Moreover, the above arguments also apply to the transmission interface 11 shown in Figure 4 of Hamstra. In particular, transmission of outgoing binary information from such a transmission interface 11 is controlled by the pointer controller 20. More specifically, and as described above for Figure 3E of Hamstra, the read pointer 30 is advanced as committed binary information 50 committed for transmission is read out of the RAM (cf. Column 10 lines 63-67). Thus, it is clear from Hamstra itself that transmission from the transmission interface 11 by reading data from the RAM 14 thereof, is the same as reading data from the RAM of the reception interface 10 of Figure 1. In neither case is a "committed read pointer" which follows a read pointer as required by independent claim 1 disclosed or suggested by Hamstra.

Accordingly, withdrawal of this rejection is respectfully requested.

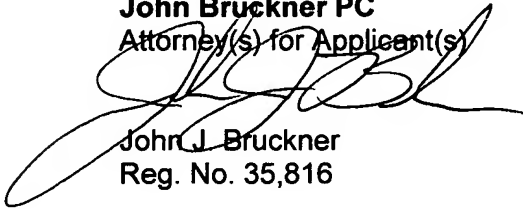
Other than as explicitly set forth above, this reply does not include acquiescence to statements in the Office Action. In view of the above, all the claims are considered patentable and allowance of all the claims is respectfully requested. The Examiner is invited to telephone the undersigned (at direct line 512-694-9145) for prompt action in the event any issues remain that prevent the allowance of any pending claims.

In accordance with 37 CFR 1.136(a) pertaining to patent application processing fees, Applicant requests an extension of time from August 21, 2004 to September 21, 2004 in which to respond to the Office Action dated May 21, 2004. A notification of extension of time is filed herewith.

The Director of the U.S. Patent and Trademark Office is hereby authorized to charge any fees or credit any overpayments to Deposit Account No. 50-3204 of John Bruckner PC.

Respectfully submitted,

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Address	Lock bit		High byte	Low byte
QID	0	RmW	Blocks to remove	Lock control
QID + 2			---	Block size (12-bits)
QID + 4		-	Tail of blocks	
QID + 6		RmW	Used blocks	Empty blocks (/Used Threshold)
QID + 8		1	Type (Single = -1/Element = RID)	
QID + 10		1	Committed Tail Block	
QID + 12		1	---	Committed Tail Offset(12-bits)
QID + 14		2	Latest Read Pointer	
QID + 16		2	---	Current Head Offset(12-bits)
QID + 18		2	Committed Read Pointer	
QID + 20		2	---	Committed Head Offset(12-bits)
QID + 22		2	---	Blocks to release

Figure 2

IN THE DRAWINGS:

Please insert the attached Replacement Sheet(s) in this Application.